How to convince ourselves that we are NOT stupid...

Prerona Challerjee STCS @TIFR

Connecting Cities









Volunleers





CHO USB Cables

Every pair of cilies should be connected



CHO USB Cables

Every pair of cities should be connected Are 5 cables enough?





7 cables are required

NO

7 cables are required

But why?











Pro con

No. of groups @ the beginning = nNo. of groups in the end = 1

Proc

No. of groups @ the beginning = n No. of groups in the end = 1 At each step the no. of groups reduce by at most 1

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So, the no. of steps read. $\geq n-1$



| 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 |
|----|----|----|---|----|----|----|----|----|----|
| 61 | 73 | 38 | 2 | 54 | 89 | 92 | 44 | 13 | 29 |

| 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 |
|----|----|----|---|----|----|----|----|----|----|
| 61 | 73 | 38 | 2 | 54 | 89 | 92 | 44 | 13 | 29 |















| 1 | 2 | 3 |
|----|----|----|
| 61 | 73 | 38 |

(1, 2, 3)















| 1 | 2 | 3 |
|----|----|----|
| 61 | 73 | 38 |





| 1 | 3 | 2 |
|----|----|----|
| 61 | 38 | 73 |





| 3 | 1 | 2 |
|----|----|----|
| 38 | 61 | 73 |





| 3 | 1 | 2 |
|----|----|----|
| 38 | 61 | 73 |

comparisons = 2 + 1 = 3





In general, the running time is:

$(n-1) + (n-2) + \dots + 1 = \frac{n(n-1)}{2} = \Theta(n^2)$

Algorithms Bubble Sort Selection Sort Insertion Sort Merge Sort Quick Sort Heap Sort

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Running Time

 $O(n^2)$ $O(n^2)$ $O(n^2)$ $O(n \log n)$ $O(n \log n)$ $O(n \log n)$
can we do any beller?

Can we do any beller?

• Nothing is known about the kind of data

Only comparisons allowed

Can we do any beller?

· Nothing is known about the kind of data

. Only comparisons allowed

Nothing better is possible

Can we do any beller?

· Nothing is known about the kind of data

. Only comparisons allowed

Nothing better is possible But why?

A sorting algorithm



Binary tree

A sorking algorithm



Depth



Binary tree

No. of comparisons

A sorting algorithm







Binary tree

No. of comparisons

Leaf of the binary tree



Final answer



Correct permutation

Proc

Total number of correct permutations possible = n!

Proch

Total number of correct permutations possible = n! \leq the no. of leaves in the tree

Proc

Total number of correct permutations possible = n! \leq the no. of leaves in the tree

The no. of comparisons required = the depth of the tree

Proof

Total number of correct permutations possible = n! ≤ the no. of leaves in the tree

The no. of comparisons required = the depth of the tree

and this is AT LEAST Log $(n!) = \Theta(n \log n)$

Matrix Multiplication

How many multiplication operations are required?

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How many multiplication operations are required?

What is currently known: The trivial algorithm takes $O(n^3)$ Strassen showed it is $O(n^{\log_2 7}) = O(n^{2.807})$ The current best algorithm takes $O(n^{2.3728639})$

How many multiplication operations are required?

What is currently known:

The current best upper bound = $O(n^{2.3728639})$

The current best lower bound = $O(n^2)$

Computability & Complexity Theory

Problem



Easy?

Hard?

Y

Decision Problems

L : Set of Strings x : Input string Question : Does x belong to L?

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Computability:

Is there an algorithm to solve this?

Complexily Theory:

If it does have an algorithm, then how much resource does it require?

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Resources we are interested in: 1. Time required 2. Space required

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P: Polynomial time algorithms



Input Length : n

P: Polynomial time algorithms Run time: $O(n^k)$ for some const. k



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NP: Non-deterministic Polynomial Lime algorithms



Input Length : n

P: Polynomial time algorithms Run time: $O(n^k)$ for some const. k

NP: Non-deterministic Polynomial time algorithms

Verifying time: $O(n^k)$ with $O(n^k)$ bits of advice for some const. k





Alan Turing









Amazing first book





MICHAEL SIPSER



Algebraic Problems



Easy?

Hard?

Algebraic Problems?

Algebraic Problems 1. Primality Test

Algebraic Problems 1. Primality Test






Algebraic Problems 1. Primality Test 2. Matrix Multiplication

Algebraic Problems 1. Primality Test 2. Matrix Multiplication 3. Integer Factorisation

Algebraic Problems 1. Primality Test 2. Matrix Multiplication x 3. Integer Factorisation 4. Algebraic Circuit Complexity



Algebraic Problems 1. Primaliky Test 2. Matrix Multiplication 3. Integer Factorisation 4. Algebraic Circuit Complexity

and many more....