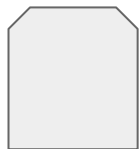
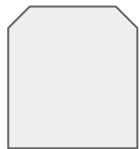
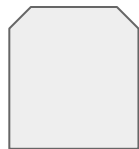
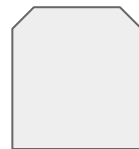
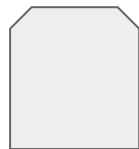
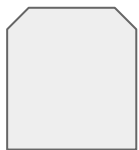
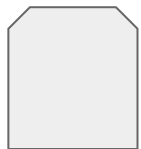


HOW HARD IS IT TO SOLVE THIS?

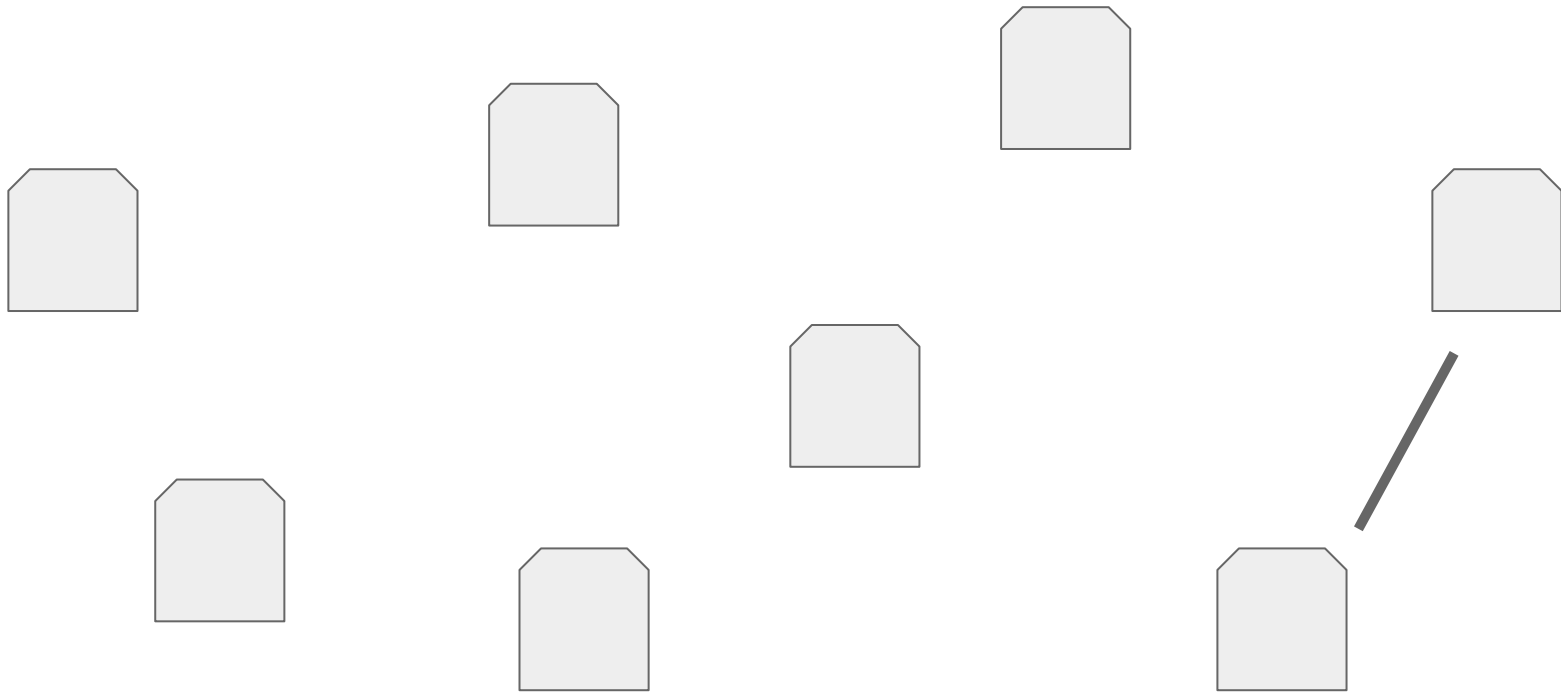
**Prerona Chatterjee**

**NISER Bhubaneswar**

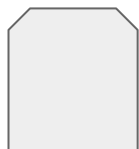
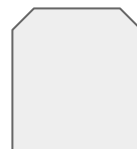
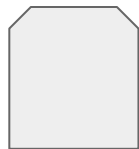
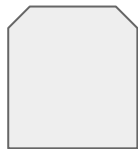
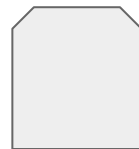
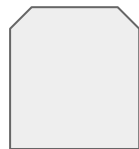
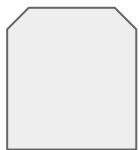
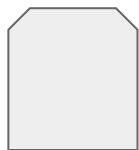
# CONNECTING VILLAGES



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- If yes, how?
- If no, why not? How much budget would be enough?

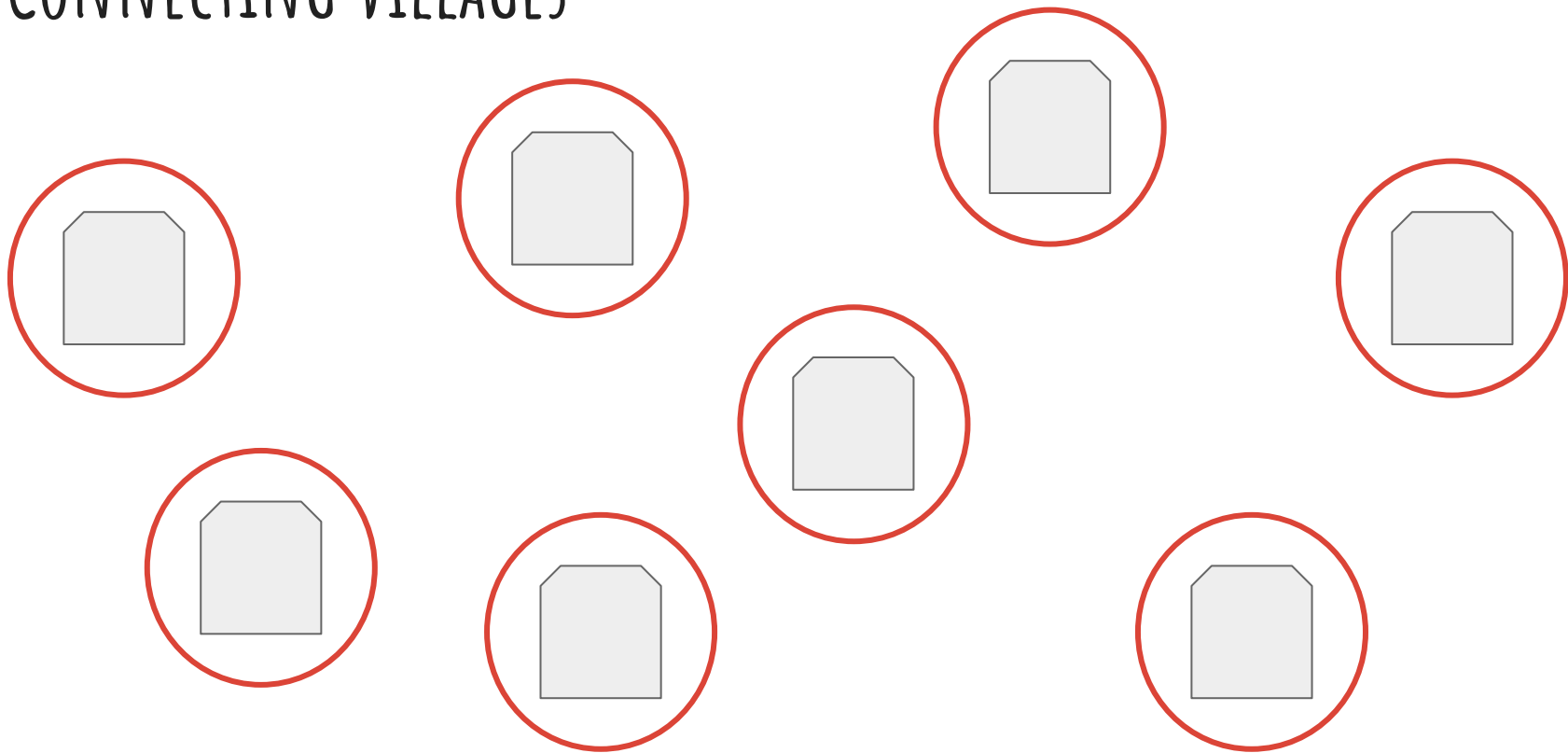
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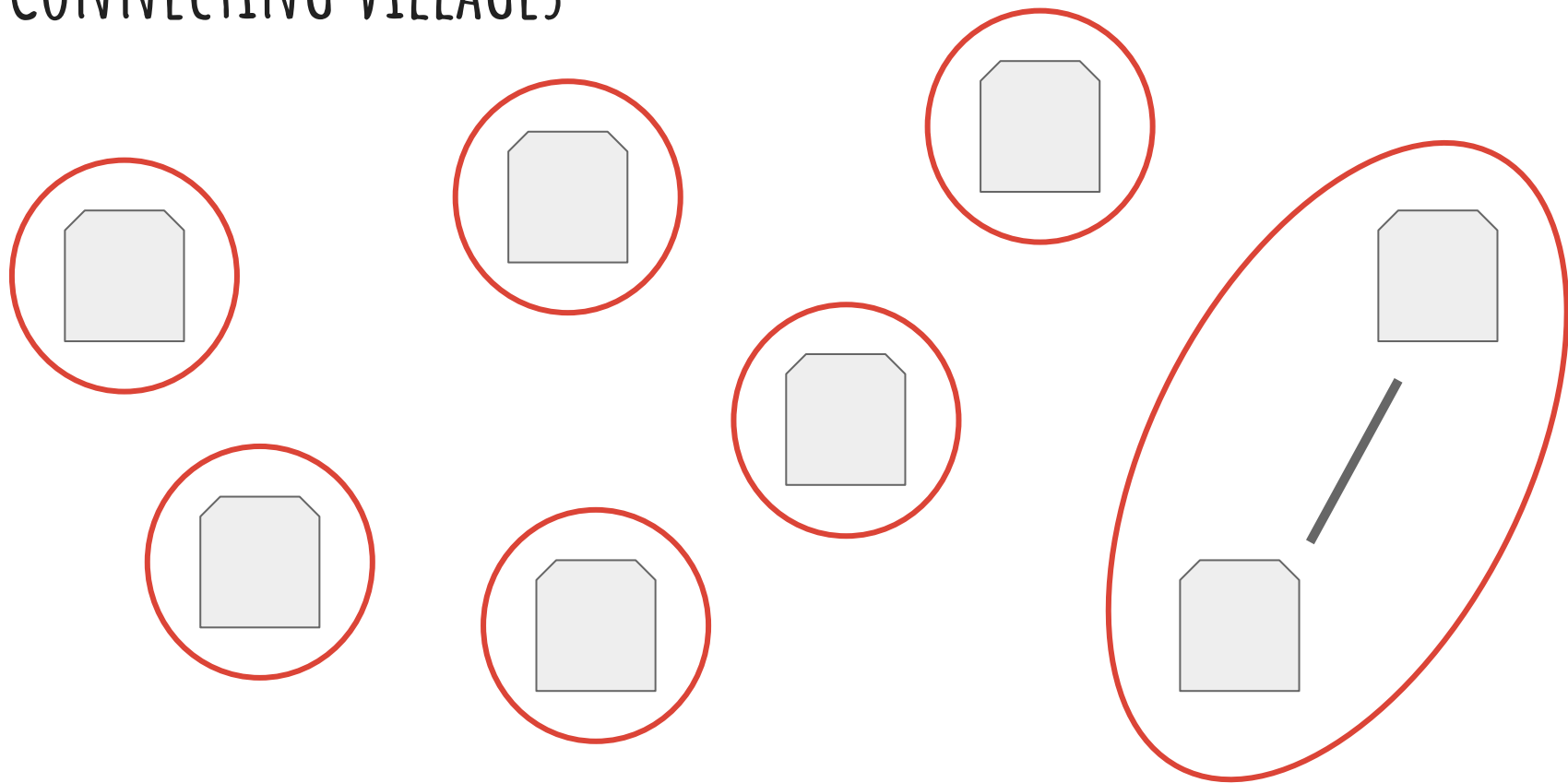
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**₹70,000 is enough. But is it necessary?**

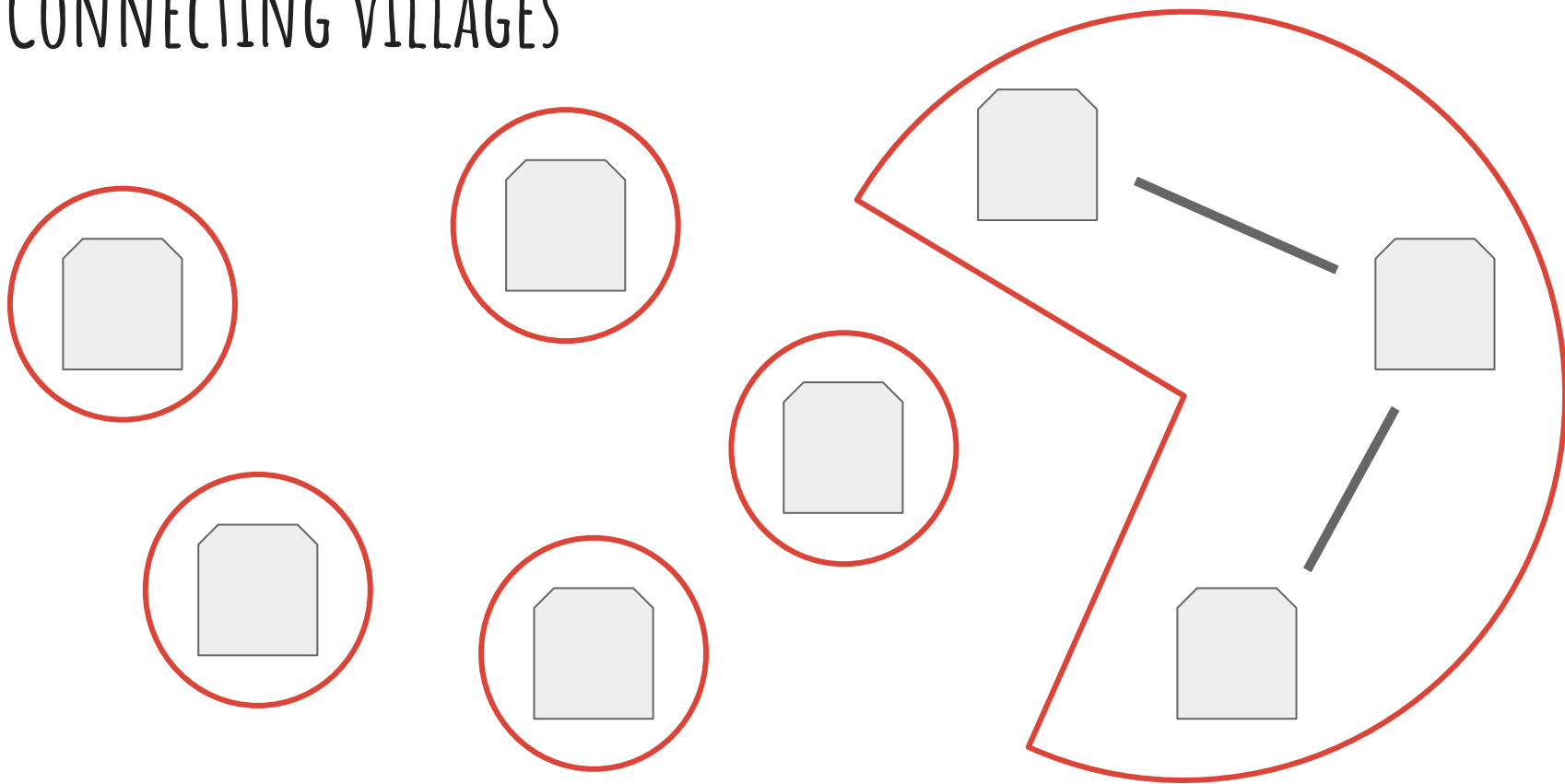
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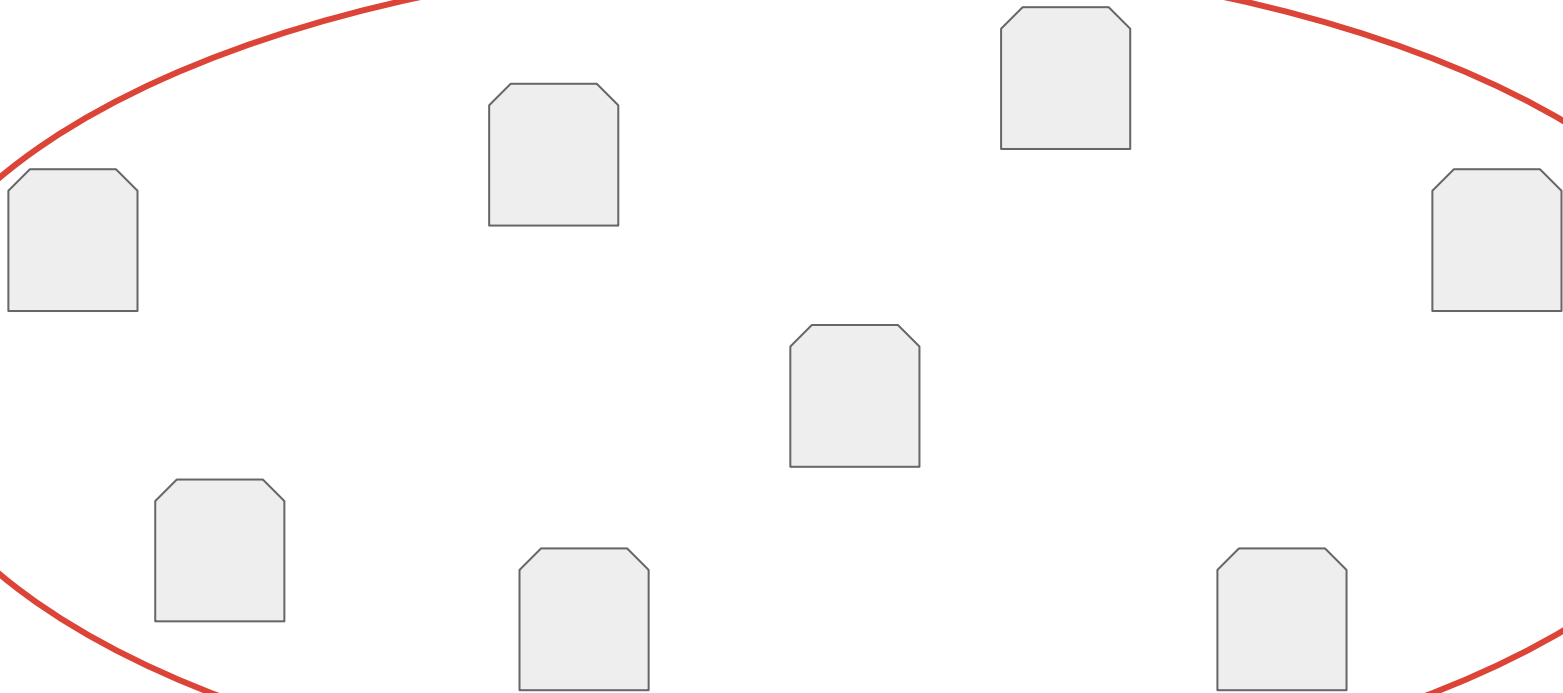
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**₹70,000 is indeed necessary.**

HOW IS PROVING  
THAT SOMETHING IS  
HARD HELPFUL?

# KNOWING THE LIMITATIONS

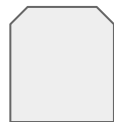
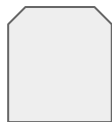
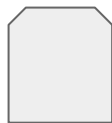
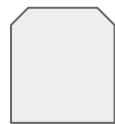
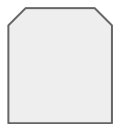
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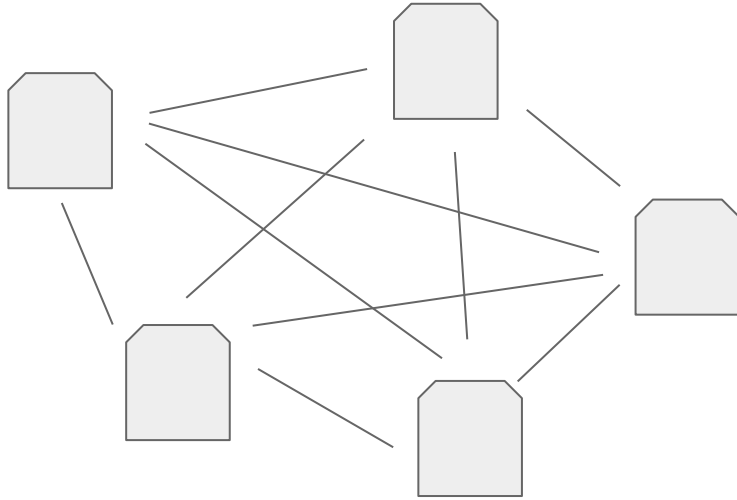
- Is it necessary to find the exact solution?

# THE TRAVELLING SALESPERSON



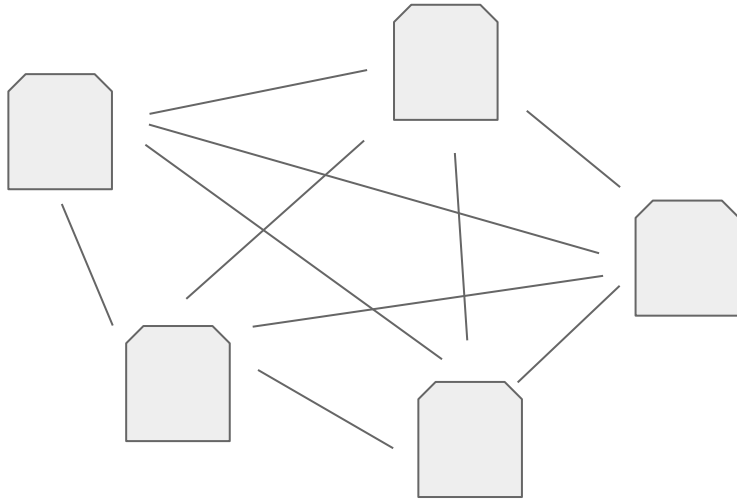


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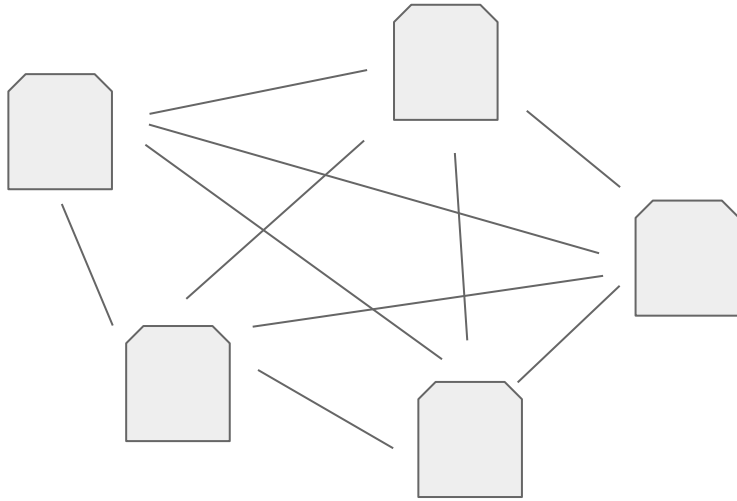
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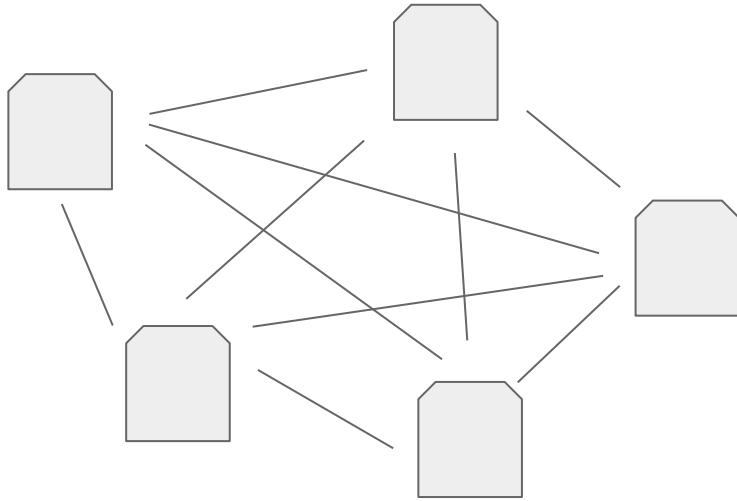
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- Different (direct) roads are of different lengths.
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- Find the route which minimises the distance.

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# APPROXIMATION ALGORITHMS

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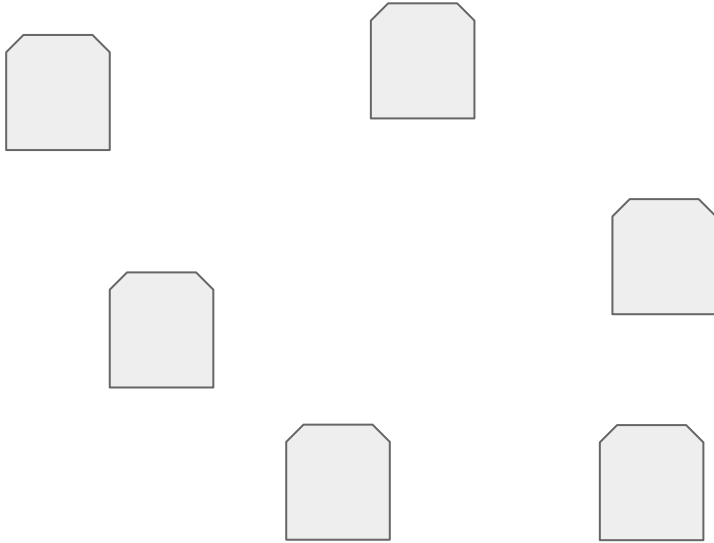
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# KNOWING THE LIMITATIONS

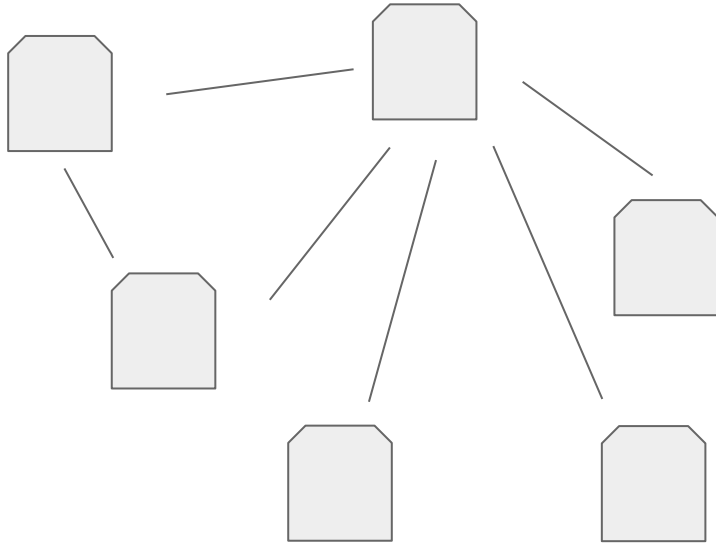
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# K-VERTEX COVER

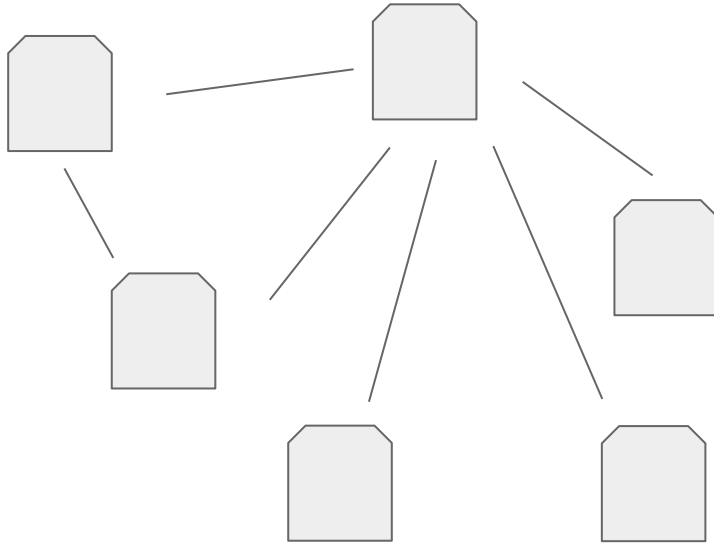


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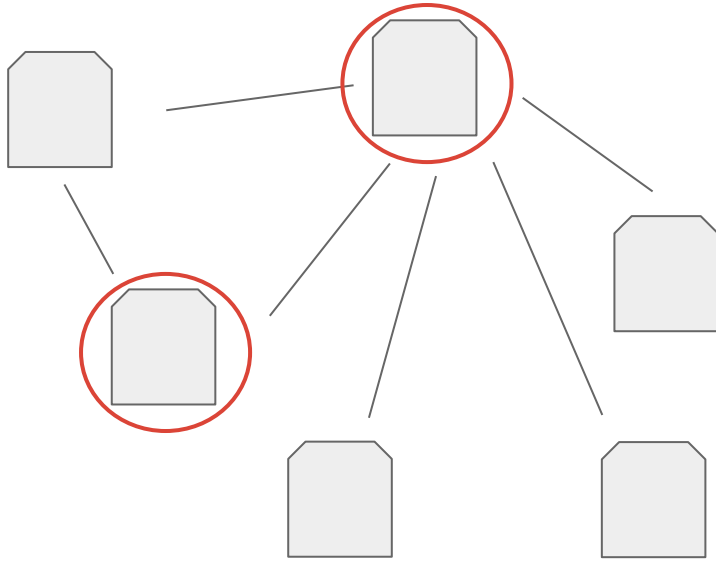
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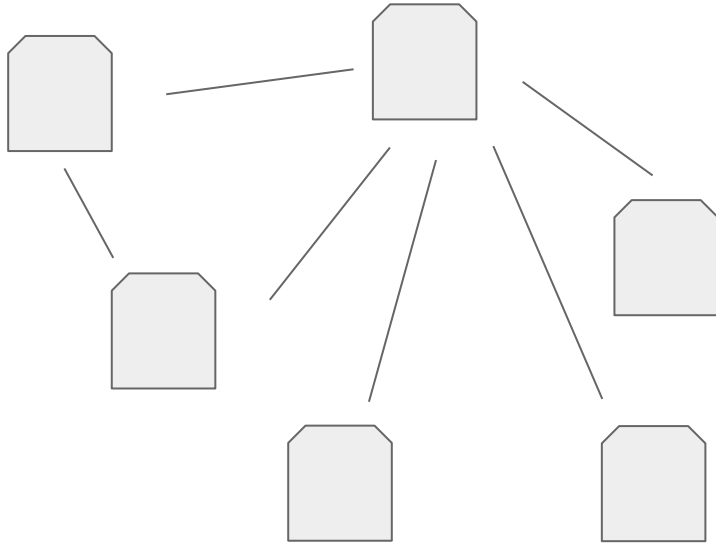
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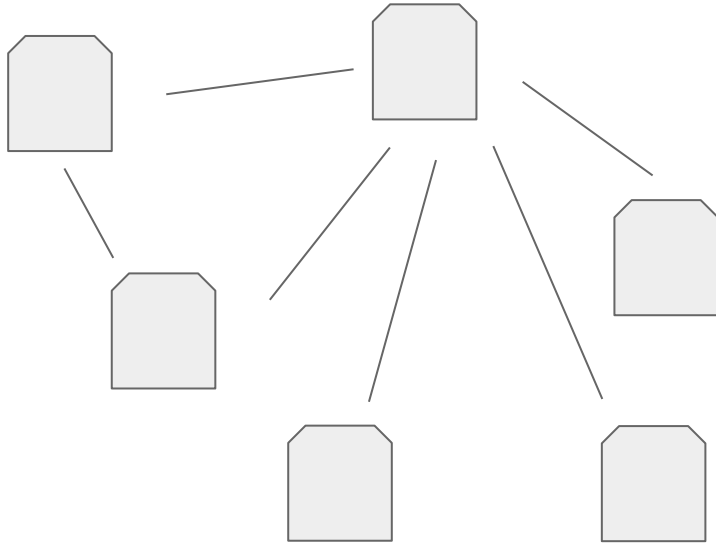
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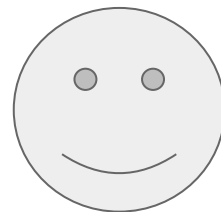
**As a group, you want to find the sum of the numbers selected by your members, WITHOUT ANY OF YOU FINDING OUT THE NUMBERS SELECTED BY THE OTHERS.**

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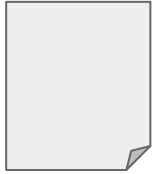
# SENDING SECRET MESSAGES



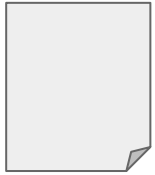
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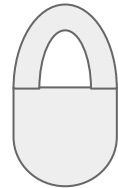
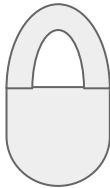
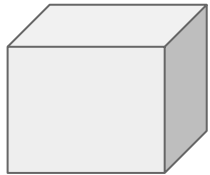
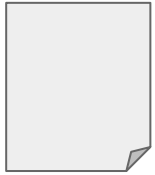
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**This protocol is secure assuming B does not have any equipment to break the lock.**



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$$G : \text{FUNC. S.T. } G(X, F(Y)) = G(Y, F(X))$$

THEORETICAL  
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STUDY THE AMOUNT OF RESOURCES NEEDED BY THE MODEL TO SOLVE THE PROBLEM.

I AM MORE INTERESTED IN  
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- Finding factors of a given number.
- Finding the shortest path between two vertices in a graph.
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- Communication models.
  - two-party/multi-party
  - compromised channels/uncompromised channels
  - broadcast/point-to-point
- Computers (a.k.a Turing machines)
- Quantum Computers.
- Circuits.
- Proof Systems.
- Prover-Verifier Games.
- .. ..

ALGEBRAIC  
COMPLEXITY THEORY

# UNDERSTANDING THE COMPLEXITY OF ALGEBRAIC PROBLEMS

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- Finding factors of a given number/polynomial.
- Computing the determinant of a matrix.
- Multiplying two numbers/polynomials/matrices.
- Finding out if a given set of vectors are linearly independent.
- Finding an annihilator for a given set of polynomials.
- .. ..



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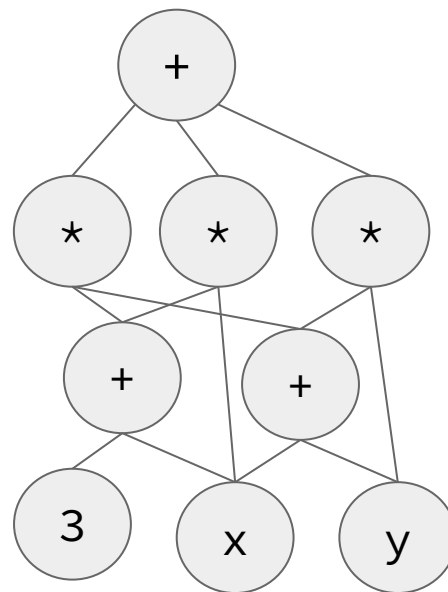
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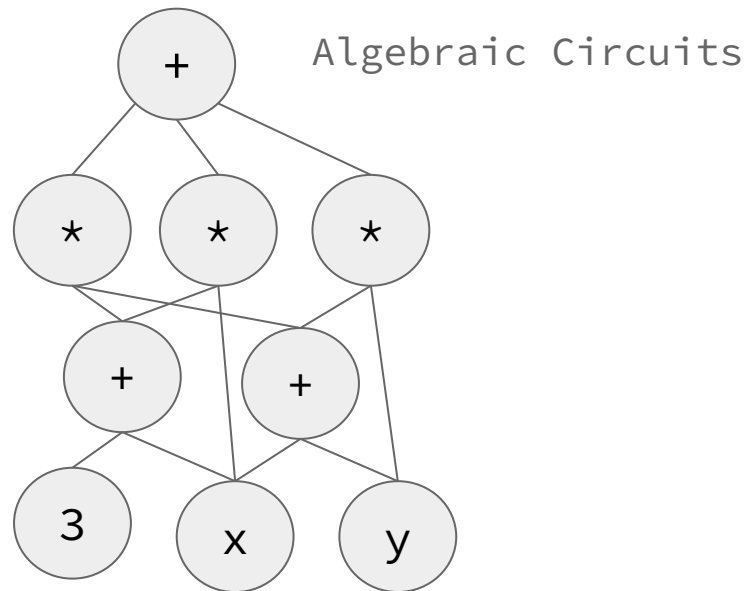
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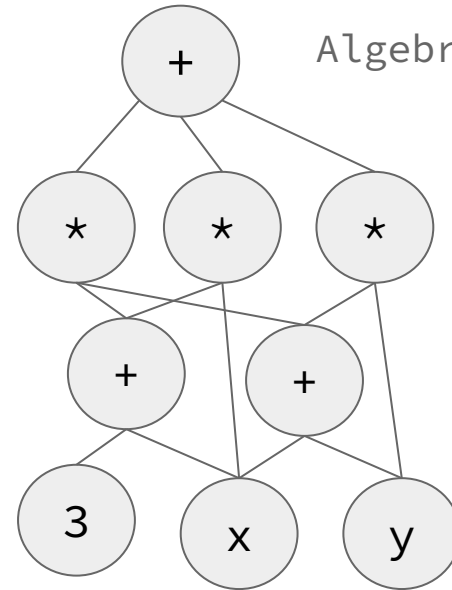
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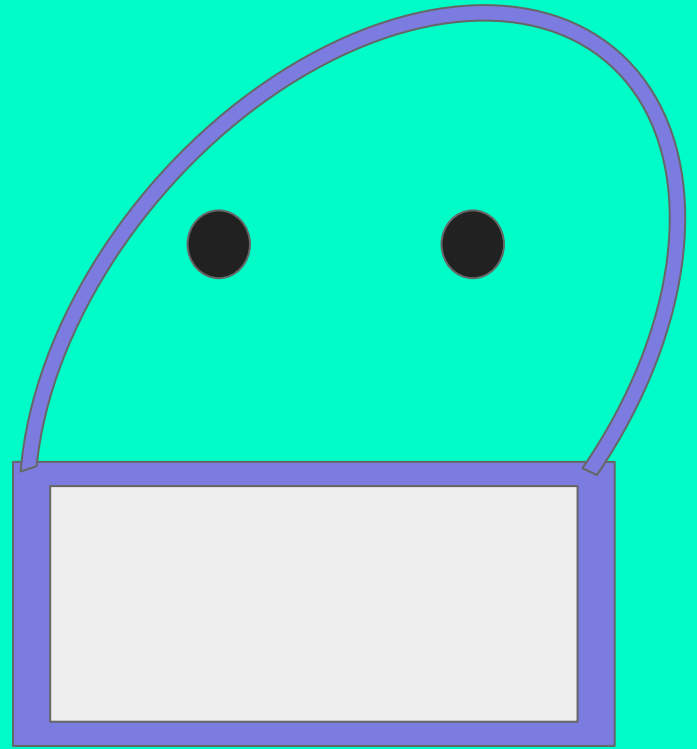
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Algebraic Circuits

- Algebraic Formulas
- Algebraic Branching Programs
- BSS Model

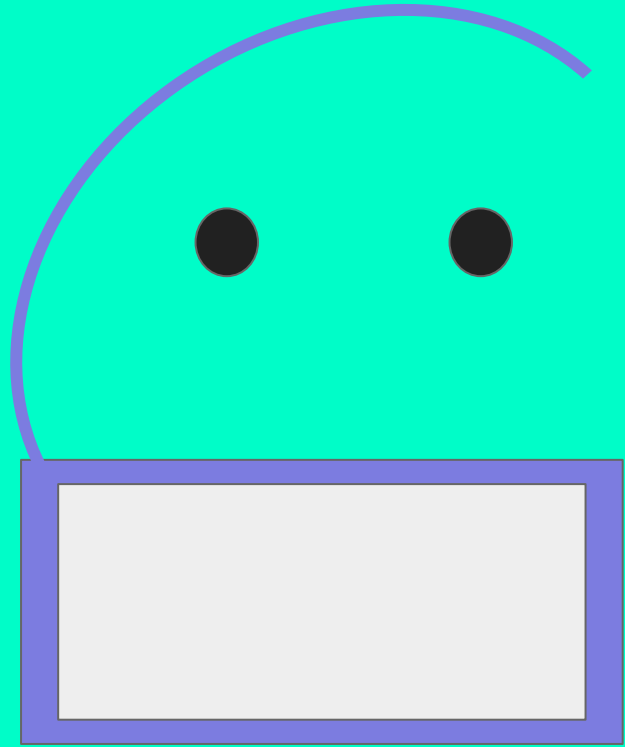
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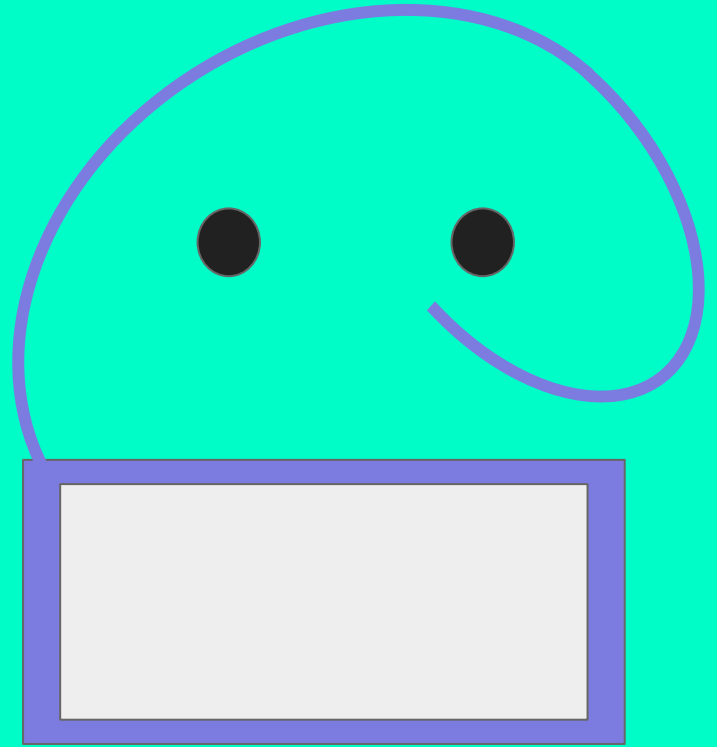
Solution: **A**



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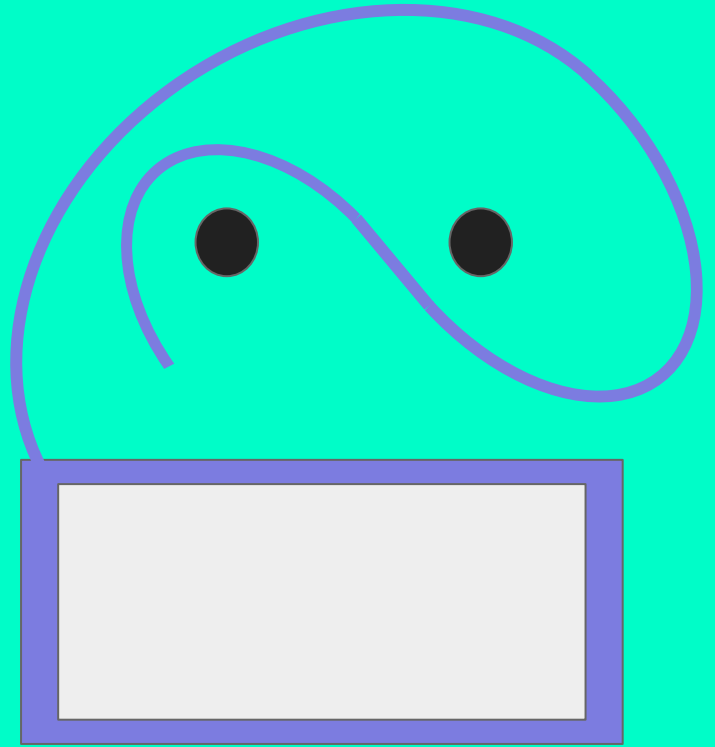
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Solution: **AB**



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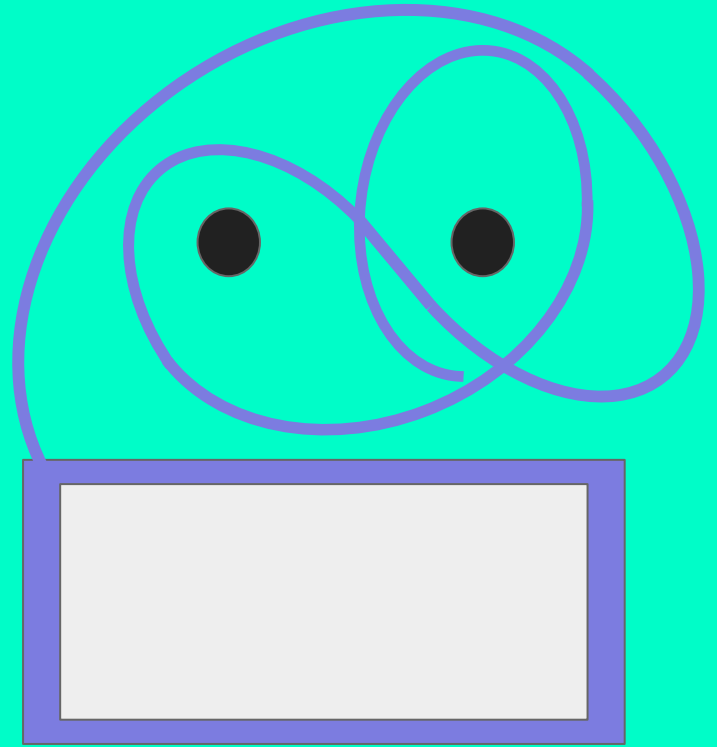
Solution:  $ABA^{-1}$





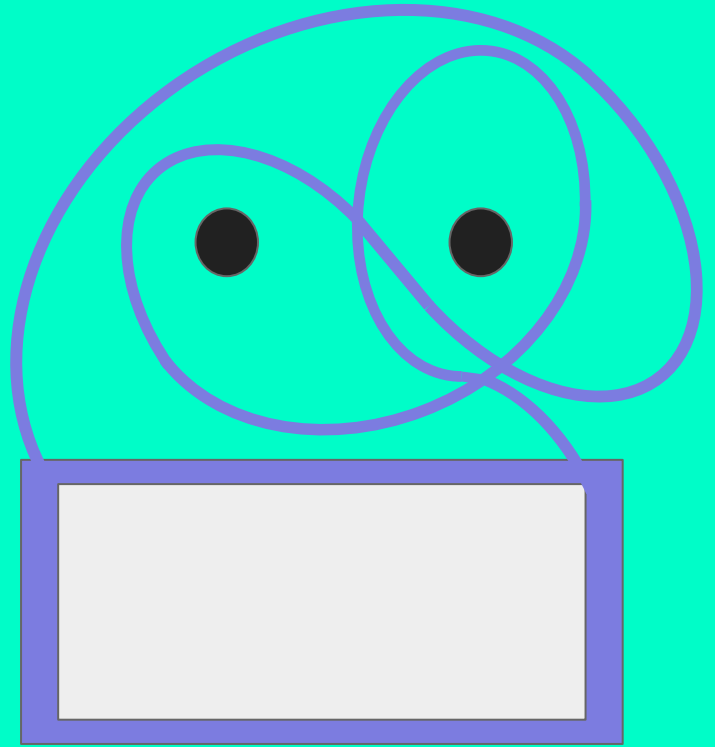
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Solution:  $ABA^{-1}B^{-1}$



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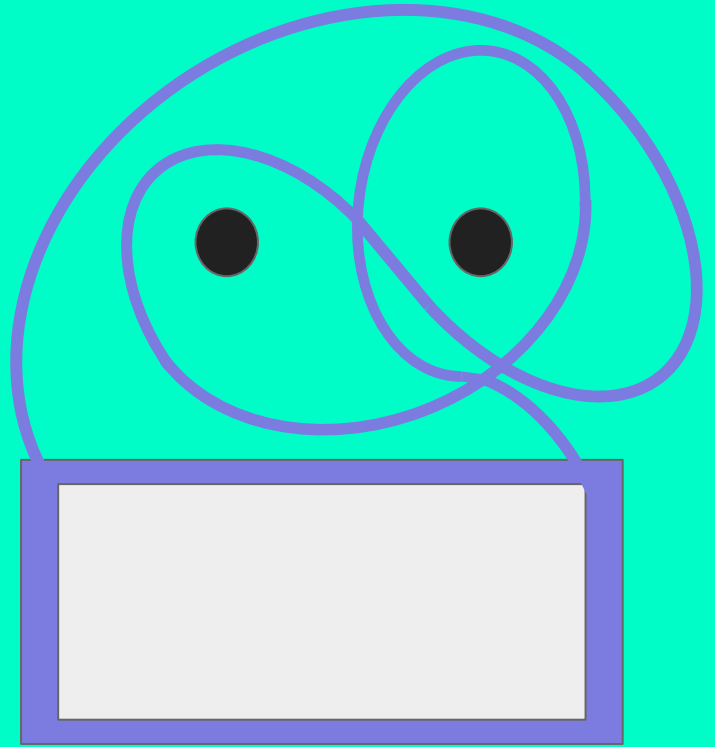
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**COMMUTATORS**



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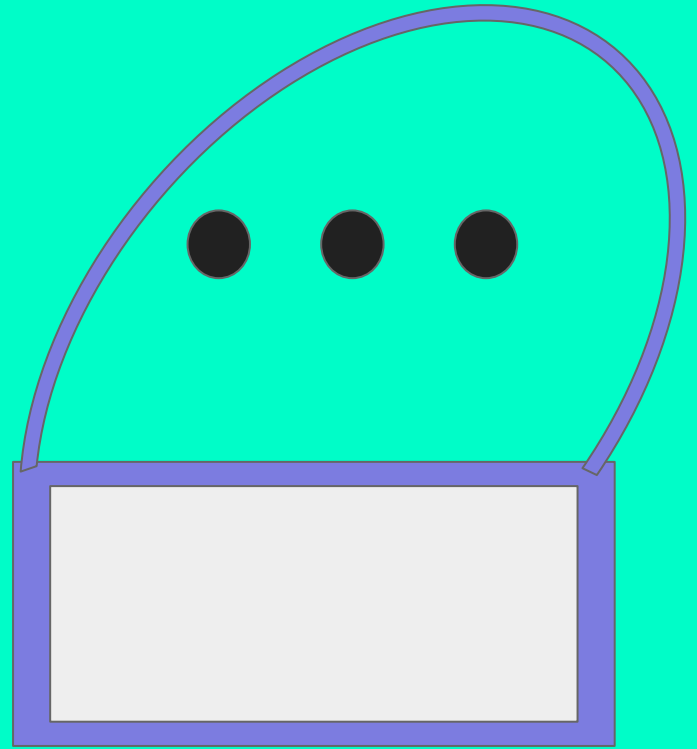
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Check out: Chai and Why?(Dec 20, 2020) by Ramprasad Saptharishi

“Commutators! Hanging pictures and solving Rubik's Cube”

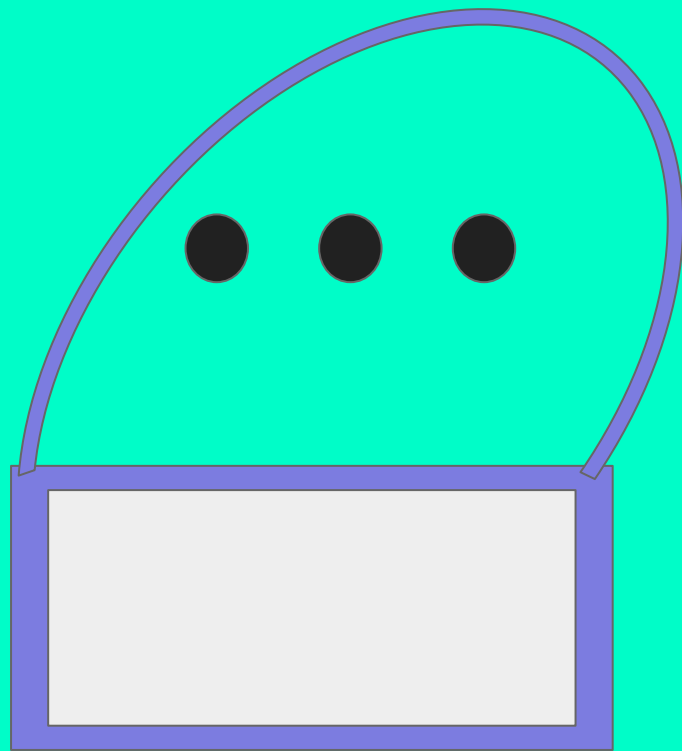
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Solution

$$A(BCB^{-1}C^{-1})A^{-1}(CBC^{-1}B^{-1})$$



SCOS @ NISER

## Secure Multiparty Computation

Secure multiparty computation (MPC) is a cryptographic protocol that **allows multiple parties to jointly compute a function** of their inputs while revealing as little as possible about those inputs.

Some of the applications of MPC are **online auctions, voting** etc.



## Data Clustering

Data clustering is the process of **grouping data points together based on their similarities**. Clustering is an unsupervised learning technique.

Application of data clustering includes **market segmentation, image segmentation, anomaly detection** and many more.

## Algorithm Design

Algorithm Design refers to developing efficient algorithms to **solve computational problems and analysing their complexity**.

The aim is to devise algorithms that optimize time and space complexities, addressing challenges in various domains like **optimization, cryptography, artificial intelligence**, and more.

## Machine Learning

Machine Learning (ML) is a type of artificial intelligence (AI) that allows software applications to become more accurate in **predicting outcomes without being explicitly programmed to do so**.

Application of Machine learning includes **healthcare, natural language processing, recommendation systems**.

## Complexity Theory

Complexity Theory is the study of different **computational models**.

Research in this area focuses on understanding the **power and limitations** of objects that model **real-world machines** with varied restrictions. Knowing the limitations of a model helps us devise **secure protocols** against them.

THANK YOU !

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<https://preronac.bitbucket.io/>